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Discovering Valuable Patterns through InternetWeb-Log Access Analysis

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ABSTRACT

This study outlines the usage of data mining techniques to analyze web log files. It shows the direction in which web mining can be performed to unearth concealed information in huge access log data. An attempt has been made to give an overview on how to derive association rules from web server data mining. Paper also discusses implementation of OLAP technology to perform web usage analysis.

Keywords. Web mining, web log analysis, usage analysis, association rules, OLAP cubes

1.0 INTRODUCTION

It is difficult to visualize the present market situation without taking World Wide Web (WWW) into consideration. Everyday new complexities are being faced in web site designing and site navigations. When the market is moving into internet, it is also a major concern for business industries to observe customer's interests. All these problems together seek for a richcontent and easily accessible site. To provide a better web designing, navigation through the web pages becomes an important input for analytical purposes. Here, Web mining refers to discovery and analysis of useful information from the World Wide Web. It also focuses its attention on user accesses data. Once the interesting patterns are recognized from web access log data, it will help the company in restructuring and better management of the web site, giving more effectiveness to it. Important point is that the web servers register a web log entry for every single access they get in which they save the URL requested, the IP address from which the request originated, and a timestamp, and with the rapid progress of WWW technology, and the ever growing popularity of the WWW, a huge number of Web access log records are being collected [3]. Frequently visited sites would easily end up with repository of hundreds of megabytes of log data. Here comes the issue of data mining considering colossal files of raw web log data where retrieving significant and useful information is a nontrivial task [3].

This study is an attempt to analyze web logs to dig information about session identification, user navigations, web usage analysis, and association rules between various pages. In order to show data mining techniques with examples, *a hypothetical company* (www.OnlineBookStore.com) is considered in this paper. The behavior of the web page readers is imprinted in the web server log files. Analyzing and exploring regularities in this behavior will significantly improve system performance, enhance the quality and delivery of Internet information services to the end user, and identify population of potential customers for electronic commerce. Thus, by observing people using collections of data, data mining will bring considerable contribution to company's web site designers.

2.0 ABOUT THE COMPANY

This hypothetical company hosts a website for online book shopping (*Site map is shown in appendix*) There are some web pages to provide the information about the company itself. Different categories of books include computers, children, sports, and fictions. A site navigation tree is presented below. Different categories of books are managed in their respective

directories e.g. computer books are kept in /Computer directory. /Home.html is the root file and all other files can be reached from this root node. Category files e.g. Sports.asp is accessible from SpecialOffer.asp. Moreover, users cannot directly access two files e.g. to reach Java2.asp from Java1.asp, user will navigate through /Computer/Computer.asp.

Note:

I.Every web page name is followed by a letter (A,B,C,...) in the site map. These letters have been used for derivation of association rules.

2. Access log data are randomly generated for analytical purposes, and so are **not** the actual data.

3.0 STEPS FOLLOWED IN WEB MIINING

- · Convert server log files into relational table format
- Data Preprocessing and Cleaning
- Session Identification
- Path Completion Analysis
- Transaction Identification
- Discovery of Association Rules
- OLAP Technology Implementation
- · Characterization and Comparison (Web usage analysis)

ACCESS LOG FILE in text format:

130.85.253.114 — [05/Apr/2002:11:34:55 +0100] "GET /Home.html HTTP/1.0" 200 2048 130.85.253.114 — [05/Apr/2002:11:35:37 +0100] "GET /Category.html HTTP/1.0" 200 1536 130.85.253.114 — [05/Apr/2002:11:36:25 +0100] "GET /Computer/ Computer.asp HTTP/1.0" 200 2048

.....

The access log file is stored on server in ASCII format. Log file contains following information about each navigation:

• IP Address of the computer the request is coming from

- User ID of the user who generated the request (if assigned any)
- *Date and time* of the request
- user action (GET or POST)
- URL of requested page
- Name and version of the protocol
- Status code (or error code) of the request e.g.
- 200: successfully received
- 400: bad request
- 505: HTTP version not supported
- Size of the page in bytes

Table2: Identification of session from log files

Session Number	Session IP Address	Start Row number	End Row Number	Access Path
1	130.85.253.114	1	5	A-D-F-M-N
2	130.85.253.114	6	10	A-L-B-H-A3
3	207.46.230.220	11	13	B-H-A4
4	207.70.7.168	14	20	A-D-G-Q-T-R-W
5	209.96.148.192	21	22	C-E
6	144.92.104.37	23	33	A-E-L-F-O-P-I-A6-H-H-A2
7	170.248.128.30	34	37	D-I-A5-A7
8	199.171.55.3	38	40	E-J-K
9	198.82.162.11	41	45	A-D-I-A7-F

3.1 Step 1 - Convert The Raw Data From Access Log Files Into Relational Data Format.

A script is run to extract and put ASCII data into respective columns in a table. Table 1 (appendix) presents web usage data restructured in relational format. (*IE denotes Internet Explorer*)

3.2 Step 2 - Data Preprocessing and Cleaning

Data preprocessing is performed on relational data. At this stage, preprocessing and cleaning processes have been discussed taking discovery of association rules into consideration. Step 8 continues further with discussion on data preprocessing for OLAP implementation.

a) Data reduction: Remove undesired fields

Data set is reduced by avoiding some attributes which do not directly influence the web mining results. Hence, from current log data, <u>attributes</u> User ID, User Action, and File Size are omitted.

b) Data cleaning:

(i) HTTP protocol requires individual connection for every file. Hence even request for an ASP page may lead to several log entries on graphics/ scripts written in log file. Concentrating on html/asp pages, records having file path ending with gif, jpeg, jpg, bmp are omitted. Here <u>row number</u> 5,20,22,24,33,43,45, 46 and 54 are removed.

(ii) Only successful requests are filtered.

If status code=200 → request successful; record is retained.

If status code!=200 → request unsuccessful; record is removed.

Here *row number 24, 37, and 56 are removed from the log data table.* These records are stored in "**error_info**" table to keep track of various errors occurring in web pages.

(iii) After removing error records, *status code and HTTP protocol fields are discarded requiring no attention for association rules derivation.* Table 2 (appendix) represents data after cleaning and preprocessing.

3.3 Step 3 - Session Identification

Discovery of association analysis from web log data starts with *transaction identification*.

Transaction can be identified by "a collection of user clicks to a single Web server during a user session". A criterion for identifying server session is implemented by examining if client has not surfed through the site for a reasonable long time.

<u>Method used to identify unique user sessions</u> If the time between page requests exceeds a certain limit, it is assumed that the user is starting a new session.

Let's assume that the **Timeout** (the session becomes inactive if not used for this specific time) is *30 minutes*.

Each log record *l* will contain this information: *l.ip*: client IP address *l.uid*: user id of the client *l.url*: URL accessed *l.time*: access time

Session can be identified by comparing all the tuples as described :

 $S = \langle ip_t, uid_t, \{(l_1^t.url, l_1^t.time), ..., (l_m^t.url, l_m^t.time)\} >$

where, for $.1 \le k \le (m-1), l'_k \in L, l'_k$ ip $= ip_t, l'_k$ uid $= uid_t, l'_{k+1}$ ip $= ip_t, l'_{k+1}$ uid $= uid_t$ and $, (l'_k time - l'_{k+1} time) \le W (= 30Minutes)$

By looking at table 2, we find that all IP addresses except 130.85.253.114 have unique user sessions. IP address 130.85.252.114 has two user sessions because time difference between row number five and six is (05/Apr/2002:13:14:37-05/Apr/2002:11:42:23) greater than 30 minutes. Hence 130.85.253.114 has two user sessions, one from row number 1 to 5, and another from row number 6 to 10. A summary **user session table** is shown below:

3.4 Step 4 - Path Completion Analysis

Even after unique user sessions are successfully identified, another problem remains if there are important accesses that are not recorded in the access log. We refer this problem as *path completion*. An approach is used here to identify missing web pages from the log file. *This approach is explained as follows:*

If a page request is not directly linked to the last requested page, it is assumed that the page is already in user's recent request history, and so user backtracked with the "back" button available on browser. This in effect calls cached version of the pages until a new page was requested. Site map can be used here to identify the missing pages. If more than one page in the user's history contains a link to the requested page, assumption follows that the page closest to the previously requested page is the source of the new request. Missing page references inferred through this approach are added to the user session file.

Session number 1: Access path is A-D-F-M-N. There is no direct path from M to N. Hence, the user must have backtracked to reach N i.e. user would have used page F to reach page N. hence complete path would be A-D-F-M-F-N.

Session number 2: Access path is A-L-B-H-A3. Again there is no direct path to reach B from L That means user has once backtracked to reach B from L via A. Hence complete path would be A-L-A-B-H-A3.

Session number 3: Access path is B-H-A4. It does not require any further modification in the access path.

Session number 4: Access path is A-D-G-Q-T-R-W. There is no direct path from R to W, so user must have backtracked to G first and then has navigated to W. hence complete access path is A-D-G-Q-T-Q-G-R-V.

Session number 5: Access path is C-E. This is a complete path considering the fact that the site has common header with links to all first level pages.

Session number 6: Access path is A-E-L-F-O-P-I-A6-H-H-A2. There is no direct path from L to F suggesting backtracking. Navigation from A6 to H suggests I as the middle page. Page H appears twice indicating removal of one page reference from access path. Hence final access path would be A-E-L-E-A-F-O-F-P-F-A-I-A6-I-A-H-A2.

Session number 7: Access path is D-I-A5-A7. Final path is D-I-A5-I-A7.

Session number 8: Access path is E-J-K. Final path is E-J-E-K. Session number 9: Access path is A-D-I-A7-F. To navigation from A7

Ses	sion Number	Session IF	Address	Start Row number
End Rov	v Number			Access Path
1	130.85.253.114	1	5	A-D-F-M-F-N
2	130.85.253.114	6	10	A-L-A-B-H-A3
3	207.46.230.220	11	13	B-H-A4
4	207.70.7.168	14	20	A-D-G-Q-T-Q-G-R-W
5	209.96.148.192	21	22	C-E
6	144.92.104.37	23	33	A-E-L-E-A-F-O-F-P-F-A-I-A6-I-A-H-A2
7	170.248.128.30	34	37	D-I-A5-I-A7
8	199.171.55.3	38	40	E-J-E-K
9	198.82.162.11	41	45	A-D-I-A7-I-D-F

Table 3: Identification of access path (path completion) from session file

Table 4: Transaction table from session table (using Maximal Forward reference)

Sessie	on Number	Access Path	
			Transaction
1	A-D-F-M-F-N		A-D-F-M, A-D-F-N
2	A-L-A-B-H-A3		A-L, A-B-H-A3
3	B-H-A4		B-H-A4
4	A-D-G-Q-T-Q-G-R-W		A-D-G-Q-T, A-D-G-R-W
5	C-E		C-E
6	A-E-L-E-A-F-O-F-P-F-A-I-A6-I-A-H-A2		A-E-L, A-F-O, A-F-P, A-I-A6, A-H-A2
7	D-I-A5-I-A7		D-I-A5, D-I-A7
8	E-J-E-K		E-J, E-K
9	A-D-I-A7-I-D-F		A-D-I-A7, A-D-F

to F, there are two possibilities, reaching either from page A or page D. Because page D is the closest to the requested page, the final path would be A-D-I-A7-I-D-F.

3.5 Step 5 - Transaction Identification by Maximal Forward Reference

This approach divides large access paths into smaller ones to identify smaller transactions. *This technique works as follows:*

Users are apt to travel objects back and forth in accordance with the links and icons provided. This senses possibility of two types of references: backward and forward references. A *backward* reference is the revisit of previously visited resource; on the other end, a *forward* reference is the visit *Table 5*

of a new resource in user session path. Transaction is defined as the set of pages in the path from the first page in a user session up to the page before a backward reference is made. When backward references occur, a forward reference path terminates. New transaction starts with next forward reference. This resulting forward reference path is termed as a *maximal forward reference*.

Considering session data from table 4, for session number 1, access path is A-B-F-M-F-N. According to the rule, first transaction will end at A-B-F-M, and so another transaction will be A-B-F-N. *Table 5 presents all transactions.*

Each of the transactions represents a basket and each resource an item. *Apriori algorithm* is now applied to derive association rules.

Iransact	ion ID Items		Transaction ID	Items	Transactio	on ID Item	S
100	A-D-F-M		106	A-D-G-R-W	112	A-H-A2	
101	A-D-F-N		107	C-E	113	D-I-A5	
102	A-L		108	A-E-L	114	D-I-A7	
103	A-B-H-A3		109	A-F-O	115	E-J	
104	B-H-A4		110	A-F-P	116	E-K	
105	A-D-G-Q-T		111	A-I-A6	117	A-D-I-A7	7
	-				118	A-D-F	
Total nun	the of transactions, $N = 19$						
Let's assu	me that minimum support le	vel = 15% = 2.	85 ~ 3				
	minimum con	ifidence level =	= 50%				
Itemset	Compart count	· • ·	-				
	Support count	Itemset	Support count	Itemset	Support count	Itemset	Support count
A	13	Itemset J	Support count	Itemset S	Support count 0	Itemset A2	Support count 1
A B	13 1	J K	Support count 1 1	Itemset S T	Support count 0 1	Itemset A2 A3	Support count 1 1
A B C	13 1 1	J K L	Support count 1 1 2	Itemset S T U	Support count 0 1 0	Itemset A2 A3 A4	Support count 1 1 1
A B C D	13 1 1 8	Itemset J K L M	Support count 1 1 2 1	Itemset S T U V	Support count 0 1 0 0	Itemset A2 A3 A4 A5	Support count 1 1 1 1
A B C D E	13 1 1 8 4	J K L M N	Support count 1 1 2 1 1 1	Itemset S T U V W	Support count 0 1 0 0 1	Itemset A2 A3 A4 A5 A6	Support count 1 1 1 1 1 1
A B C D E F	13 1 1 8 4 5	J K L M N O	Support count 1 1 2 1 1 1 1	Itemset S T U V W X	Support count 0 1 0 0 1 0	Itemset A2 A3 A4 A5 A6 A7	Support count 1 1 1 1 1 2
A B C D E F G	13 1 1 8 4 5 2	Itemset J K L M N O P	Support count 1 1 2 1 1 1 1 1	Itemset S T U V W X Y	Support count 0 1 0 0 1 0 0 0	Itemset A2 A3 A4 A5 A6 A7	Support count 1 1 1 1 2
A B C D E F G H	13 1 1 8 4 5 2 3	Itemset J K L M N O P O	Support count 1 1 2 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	Itemset S T U V W X Y Z	Support count 0 1 0 0 1 0 0 0 0 0 0	Itemset A2 A3 A4 A5 A6 A7	Support count 1 1 1 1 2

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3.6 Step 6 - Discovery Of Association Rules From Transaction Data

Because the **minimum support is 3**, the next picture shows frequent 1itemset **L1** and subsequent analysis.



Hence, Apriori algorithm provides these frequent item sets (resource sets) for the transaction data: $\{A\}$, $\{D\}$, $\{E\}$, $\{F\}$, $\{H\}$, $\{I\}$, $\{A, D\}$, $\{A, F\}$, $\{D, F\}$, $\{D, I\}$, $\{A, D, F\}$

<u>Strong association rules for 2-itemsets:</u> Minimum confidence = 50%

Rule 1: $\forall x \in transaction, contains(X, A) \Rightarrow contains(X, D)$

Confidence =
$$\frac{\text{Support (A, D)}}{\text{Support (A)}} = \frac{6}{13} = 46.15\%$$

Rule cannot be established.

Rule

 $\forall x \in transaction, contains(X, D) \Rightarrow contains(X, A)$

Confidence=75% i.e. this rule is a strong rule. It implies that 75% of the users navigating page D (/Category.html) also are visiting page A (/Home.html).

Rule

 $\forall x \in transaction, contains(X, A) \Rightarrow contains(X, F)$ Confidence=30.76%, rule **cannot** be established.

Rule 4: $\forall x \in transaction, contains(X, F) \Rightarrow contains(X, A)$ Confidence=80% i.e. *this rule is a strong rule*.

Rule

 $\forall x \in transaction, contains(X, D) \Rightarrow contains(X, F)$ Confidence=37.5%, rule **cannot** be established.

Rule

 $\forall x \in transaction, contains(X, F) \Rightarrow contains(X, D)$ Confidence=60%, rule is established as strong rule.

Rule

 $\forall x \in transaction, contains(X, D) \Rightarrow contains(X, I)$ Confidence=37.5%, this rule **cannot** be established

Rule

 $\forall x \in transaction, contains(X, I) \Rightarrow contains(X, D)$ Confidence=75%, association rule is strong.

Strong association rules for 3-itemsets:

Rule

Rule

Rule

 $\forall x \in transaction, contains(X, A)^{\wedge} contains(X, D) \Rightarrow contains(X, F)$

Confidence =
$$\frac{\text{Support (A, D, F)}}{\text{Support (A, D)}} = \frac{3}{6} = 50\%$$

Strong association rule is established. It infers that 50% of the time, if the client visits both /Home.html and /Category.html, he/she will also visit / Computer/Computer.asp.

 $\forall x \in transaction, contains(X, A)^{\land} contains(X, F) \Rightarrow contains(X, D)$ Confidence=75%, so this is a strong association rule.

 $\forall x \in transaction, contains(X, D)^{contains}(X, F) \Rightarrow contains(X, A)$ Confidence=100%, so %, it qualifies as strong association rule.

3.7 Step 7 - OLAP Technology Implementation

Data Preprocessing:

To construct multidimensional cube technology, we first need to convert raw data into various dimensions with some defined to facilitate generalization and specialization.

• Access time is represented by **TIME** dimension with *schema hierarchy*.

second < minute < hour < day < month < year < All

• URL stores file structure by server domain (if multiple servers running simultaneously), directory (where the file resides), file name, extension (.asp, .html, .cgi etc.). Schema hierarchy:

file extension<file<directory<server domain<All

• Client IP Address defines organization name and domain with operation-derived hierarchy.

Organization name<domain name<All

• **Time Spent is** difference in access time for the current page and the next page. A *set-grouping hierarchy* would then be:

If t = time spent (in seconds), {very_short_stay, short_stay, moderate_stay,

 $\{0 < t < 30\} \subset very_short_stay$

 $\{60 < t < 60\} \subset \text{short}_\text{stay}$

 $\{60 < t < 180\} \subset moderate_stay$

 $t>180 \subset long_stay$

• Range hierarchy on File_Size (in bytes):

{tiny, small, medium, large, huge} **all** (File Size} {0<File_size<1000} **t** tiny

 $\{1000 < File_size < 2000\} \subset small$

7:

8:

1:

2:

3:

3:

5:

6:

2:

```
{2000<File_size<4000} 	colored medium
{4000<File_size<5000} 	colored large
{File_size>5000} 	colored huge
• Set-grouping hierarchy for Type_of_Resource:
    [script, images] call (Type_Of_Resource)
    {asp, html} cscript
```

{jpg, jpeg, gif, bmp} ⊂ images

```
    Browser type
```

Browser version<Type_of_browser<all.

We can construct OLAP cubes with abovementioned dimensions. OLAP operations drill-down, roll-up, slice and dice, can be performed to view and analyze web log data from different angles. One such *measure* is "**number of hits**" at a particular level of granularity. Step 8 explains more on OLAP operations.

3.8 Step 8 - Characterization and Comparison

a) Finding "TOP N" requested URLs

Selecting top N pages is achieved by performing drill-down on URL dimension to file, and **all** to other dimensions. For this sample data, top 6 pages in decreasing order of count would be

1) /Home.html, 2) /Category.html, 3) /Sports/Sports.asp, 4) / CompanyInfo.html, and 5) /Computer/Computer.asp, and 6) /Fiction/ Fiction.asp

b) Comparison of browser types for web site access

Web mining can be performed on the browser type to find out commonly used browsers, and more information on their compatibility for the web site can be further studied. Pie chart below presents web site usage by browser type obtained by drilling to browser dimension, and "**all**" for other dimensions.



It explains that *Internet Explorer is the most frequently used browser*. Web site can be redesigned to provide maximum support for Internet Explorer and Netscape.

c) Comparison of various domains on web site access

The access log file is processed to convert the IP addresses into host names. E.g. www.umbc.edu has education domain.

It is easy to mine which kind of domains the web site is receiving the requests from. Pie chart explains that web site is mostly used by commercial sites (**37%**) and educational institutions (**37%**). It is obtained by drilling on IP address to 'domain name', and "**all**" along other dimensions.



4.0 CONCLUSION

As web is a one of the biggest repositories ever built, analyzing web access logs can help us understand the user behavior and would lead to better web site management by making use of Web Usage Mining. Furthermore, more information about user agents and referring resources can be collected to reveal better information. Several tools like NetTracker, SAS Webhound, Analog, and WebTrends Log Analyzer can be used to explore web log data in efficient way.

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4.White paper on Speed Tracer: A Web usage mining tool (http://www.research.ibm.com/journal/sj/371/wu.html)

5.Log File Basics

(http://slis-two.lis.fsu.edu/~log/basics~1.htm)

6.WWW access statistics for the last 12 months (http://www.netstore.de/stats/ www2002/frames.html)

7.Hits summary detail

(http://www.ideva.com/reports/ideva/idevahitsummary.htm)

8.Request detail report

(http://www.ideva.com/reports/ideva/idevarequests.htm)

IP Address	Host Name	Domain Type	IP Address	Host Name	Domain Type
130.85.253.114	www.umbc.edu	Education	144.92.104.37	www.wisc.edu	Education
207.46.230.220	www.microsoft.com	Commercial	170.248.128.30	www.accenture.com	Commercial
207.70.7.168	www.infotech.com	Commercial	199.171.55.3	www.sba.gov	Government
209.96.148.192	www.myvirginia.org	Organization	198.82.162.11	www.vt.edu	Education

APPENDIX Site map



10.4	CINCLE MINING	User	Axes Time	Lser	TilePath	1	Status	1.14	Stower
20 8 600	0.000-0.000-0.000	D.	State Busidestand	Sec.	ast borne -	Protocol	- 546	and .	Tpo .
•	131 88 293 111		1.0° ¹ 458051134422	C-87	/Bone Mul	11111	400	7748	Ŧ
2	100.05.200.111)/AL_20202110527	OPI	ACale yo yo Istal	1:154.0	:0)	1206	T
3	130.36.263.111)./Ag=2:02110445	CPI	AC: mayone a C < magnine a corp	H:HIA0	20)	2(48	2
τ.	138.35 253.114		JUAN 1158 12	GEI	ComputedFarit asp	TUPHIN -	200	1.52	エ
9	100.06.200.111)/Age2*02113632	OPI	ComputerParL jpg	1:154.0	:0)	+(S6	I
6	130.35.263.111)/.0.px2=0211>243	GBI	K's myntesPan2.wy	H. HEUN	20)	1792	3
1	138.35 253.114		JUATE2 02131407	GEI	/House Mand	101610	200	20,46	Sintakan V 2
\$	130.36.263.111		02/Mgm2=02136565	CPI	AL-WAY	H:HIA0	200	718	Sataspi (2
9	130 58 253 111		0.00 ¹⁴ 2005131243	C-87	Ирази Пійтлар	4-1120	101	10.34	Wetcops 6.7
	130.55.263.114)/A;#2=02151851	GEI	/ទីរចនា/ទីរួចនា រនៈទ	11154.0	200	1048	Жатақта 4 2
-11	130.36.263.111)./Ags2 02134061	CPI	1820str/Scripthallury	H:HIA0	200	1230	Shtaspi (2
- 2	201.00/2011228		JUK12202035153	GEI	/бразв. Шатар	TUPPIN	200	10.24	Upera 401
э	ZD .46 200.220)/ATES02035333	GEI	/វីរូចនា/វីរួចនារអង្	11154.0	200	2048	Opera 6.01
1	205.46.230.220).0. <u>r</u> z2=02:05>4:31	CBI	Most Bushell up	3:114.0	200	2048	Opera 4.01
-5	201.00.168		0.2Kgg2002.0F11.19	GEI	Attoms Jahnsi	1011010	200	20,46	エ
-6	205.70.5.168		07/hga2=02014343	CPI	Area of the second	H:HIAO	20)	1234	ヱ
-7	205.70.3.168		070.gzt2=02.04.04.sl3	CBI	Childmany	H:HIGO	20)	1234	3
ð	201.00.1.168		07/6722=02011737	GEI	Children A. get 6.47	11154.0	200	22.04	I
-9	205.20.5.168		07/hga2=020128H3	CPI	Children Gidl or p	3:114.0	20)	22:40	고
30	205-20-5-168		170.14702014244	OBT	Children/Eidl gf	4*117.0	101	46.56	TP I
21	201.00.1.168		07/0412240201012401	GEI	Chillen A.go_12.w1	11154.0	200	ZC 40	T
22	205.70.5.168)7.0.gz2=02.01.45.43	CBI	Children Ci B. of	3:110.0	200	4054	Э
2.3	201.00.168		J.2Kg#2020120123	GEI	Children-Kidt .scp	TUMIN	200	20,46	보
24	205.78.5.160)7/h_w2=02.01.15.22	OPI	AChillenetfielt. g.f.	1:154.0	40)	0	ニ
25	209.06.18.192)80. <u>1</u> #2=021243.49	GBI	AboutCompany.html	3:1210	20)	1792	Monilla 4.0
25	209.96. (8.192)&ATT2=02192341	C-EI	AC) mpanying .html	11154.0	200	1(2)	Mosfila \$.0
27	144.32.104.37)39kga2*021120.02	OPI	للماط مع 100	1:154.0	:0)	2(+0	1 II
38	141.42.114.38		190kg#7#071144430	C-PT	All many and the	4,123.0	101	10.24	7
29	144.32. E4.3T		090AIX2=021+2723	C-EI	100 W 101	11154.0	200	718	T
70	14 .02. E4.3)90.pr2=02114662	CBI	Computer Computer to P	J. H. O	200	2(48	38
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32	144.32.104.37		02110210	OPI	Computerly and my	1:154.0	:0)	1(2+	ヱ
33	144.02.104.38		099 52#2#021 16330	CPI	No agrate for al stat	H:HIA0	20)	4054	z
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35	144.32.104.37		0394142-0211-07-0	eri	/Luindardy.og	1.154.0	:0)	2(+0	<u> </u>
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38	144.02.104.37		00%pa2=0211+027	CPI	/Boosth/Bgo th say	3.114.0	20)	2(48	2
.9	164.52.116.30		J90KT1202115331	GEI	18post-18occall asy	1:1010	200	1048	<u> </u>
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12	110.240.120.00		13AL101071137	OPI	arranged and by	1.110	200	1034	
11	170.2(8.128.30		110.142.02071231	GHI	Beton Line op	35150.0	200	2048	Siste op 4 2
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50	190.52. 62.11		CANTER OF 020355 31	CEL		1.1010	200	11.46	1 4
51	190.52. CZ 11		LIATE2 02 0 501 X 6	OPI	A SERVICE	1.11410	200	12.56	14
52	1946-1521-1622-11		110.220202039257	GBI	NUCLOADER F.F.	2:12010	200	22.04	
30	190.52.102.11		130 2002030413	OT	arconnu rana ap	1 ININ	200	21,46	<u> </u>
34	190.32.102.11		13/K147002033630	orr	Alexander and it is	1:164.0	200	+(56	14
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- 30	130.52. CZ 11		LIATE2 01037531	GEI	ALCONG ATTAL	E 11109	1 202	Q.	1 - 2

Table 1: Access log file in tabular format

Table 2: Web log data after preprocessing and cleaning

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2	130 85 253.1.4	03///mr2002:11:3557	Category html	E
3	130 85 253 114	05/8mm200211/3625	ComputerComputers sp	. जा
4	130 85 253.114	03/4 2000:11:38:12	ComputerDerllap	Ē
5	100 05 250,114	03///mr/2002:11:42:23	ComputerPer2 as	E
0	130 85 253,1-4	05//5002:13:1437	Homerml	New Cape 6.2
7	13U 8t 253,114	03/05/07/2012 (13:15:35)	1×7540	Netscape 6.2
8	130 85 253 114	03///mr2002:13:17:13	SpecialDEeraso	Netscape 6.2
9	130 85 253 114	05/877202-13-18:21	StorsSportcasp	Netscore 6.2
10	130 85 253 114	05//13/2031	Sports Ra bethellasp	Netscape 6.2
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13	207 41 230 220	10/3003412-11:441-1	Shorrs/Rasetallasp	Uperaŭ II
14	207 70 7.168	07/811-20.02.04:11:19	Hometini	L E
16	207 70 7 168	00///mm200.04/13/23	Category html	<u> </u>
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38	199 171 55.3	1:/// 2006 20 33:0	Companying hall	1.vnx383
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